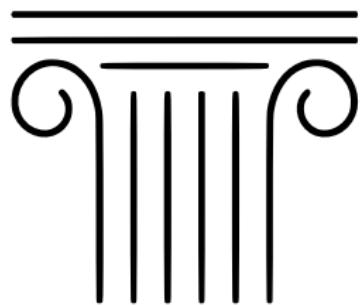


# **Games and Strategy TA 1**

*PD, Nash Examples, Oligopoly*



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## Prisoner Dilemma, example

		Player 2	
		C	D
Player 1	C	3, 3	0, 5
	D	5, 0	1, 1

**Figure:** An example of Prisoner Dilemma

- C → cooperate; D → defect
- What is the Nash equilibrium?



## Prisoner Dilemma, example

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**Figure:** An example of Prisoner Dilemma

- C → cooperate; D → defect
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## Prisoner Dilemma, general formulation

		Player 2	
		C	D
Player 1	C	A, A	B, C
	D	C, B	D, D

Figure: The general Prisoner Dilemma

A game is a PD if...

- $C > A$  and  $D > B$  i.e. defecting is always the best choice
- $A > D$ , i.e. both players would gain from cooperation

Hence, a PD features  $C > A > D > B$ .

$C$  is the *temptation*,  $B$  is the *sucker's payoff*.



## Nash equilibrium

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### Definition (Nash Equilibrium)

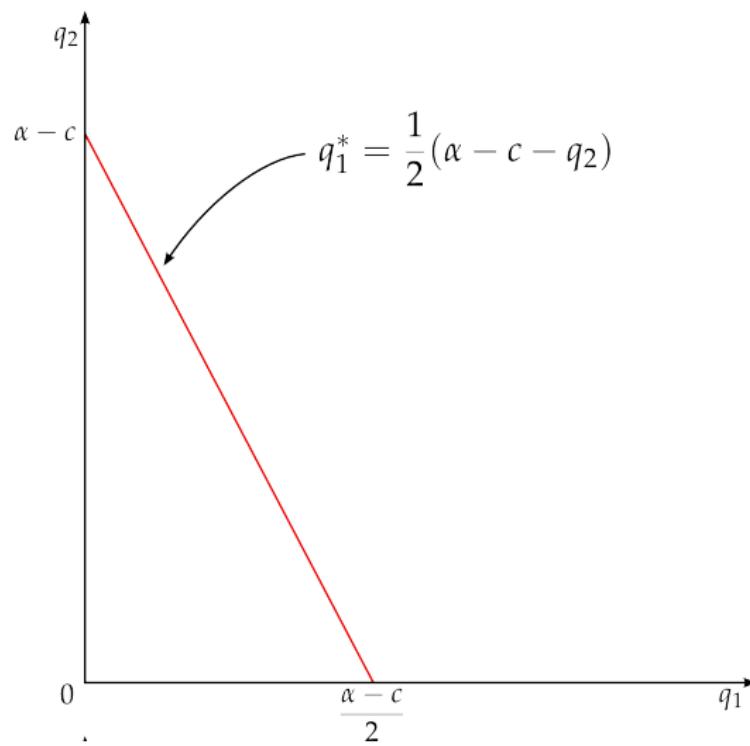
A *Nash Equilibrium* is an action profile  $a^*$  with the property that no player  $i$  can do better by choosing an action different from  $a_i^*$ , given that every other player  $j$  adheres to  $a_j^*$ . Or, more formally

$$u_i(a^*) \geq u_i(a_i, a_{-i}^*), \forall i$$

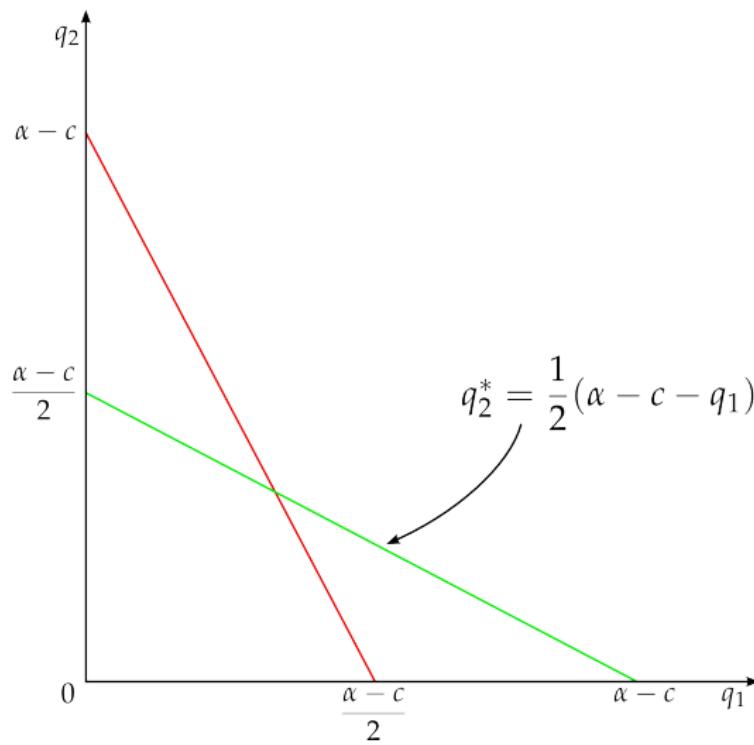
- in a Nash equilibrium, each player *best responds* to other player's actions
- which are in turn best responses.
- No agent has an incentive to deviate, to change his action from  $a_i^*$ .
- In exercises, we will heavily exploit the *no-deviation* property.



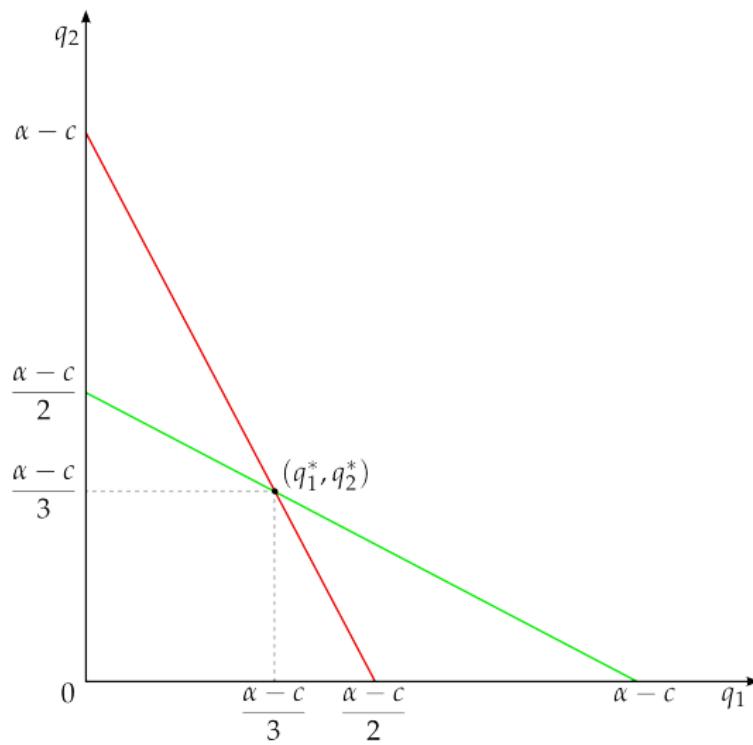
## Cournot Oligopoly - Graphic



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## Cournot Oligopoly - 2, 1, many!

	Firms	$q_i$	$Q$	$P$	$\Pi_i$
Oligopoly (C)	2	$\frac{\alpha - c}{3}$	$\frac{2(\alpha - c)}{3}$	$\alpha - \frac{2(\alpha - c)}{3}$	$\left(\frac{\alpha - c}{3}\right)^2$

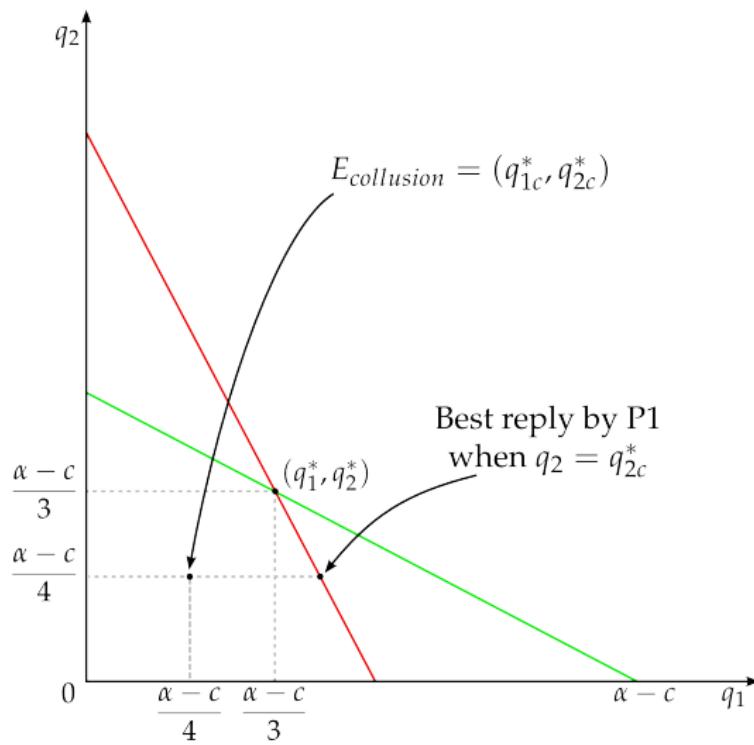


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Monopoly	1	$\frac{\alpha - c}{2}$	$\frac{\alpha - c}{2}$	$\alpha - \frac{\alpha - c}{2}$	$\left(\frac{\alpha - c}{2}\right)^2$



## Why collusion not Nash? - Graphic



## Cournot Oligopoly - 2, 1, many!

	Firms	$q_i$	$Q$	$P$	$\Pi_i$
Oligopoly (C)	2	$\frac{\alpha - c}{3}$	$\frac{2(\alpha - c)}{3}$	$\alpha - \frac{2(\alpha - c)}{3}$	$\left(\frac{\alpha - c}{3}\right)^2$
Monopoly	1	$\frac{\alpha - c}{2}$	$\frac{\alpha - c}{2}$	$\alpha - \frac{\alpha - c}{2}$	$\left(\frac{\alpha - c}{2}\right)^2$
Many	$n$	$\frac{1}{n+1}(\alpha - c)$	$\frac{n}{n+1}(\alpha - c)$	$\alpha - \frac{n}{n+1}(\alpha - c)$	$\left(\frac{\alpha - c}{n}\right)^2$
As $n \rightarrow \infty$	$\infty$	$\rightarrow 0$	$\rightarrow \alpha - c$	$\rightarrow c$	$\rightarrow 0$

